

Lecture 19 - Makeup for WrittenTest2

(\approx 90 minutes)

Lecture

Recursion: Part II (continued)

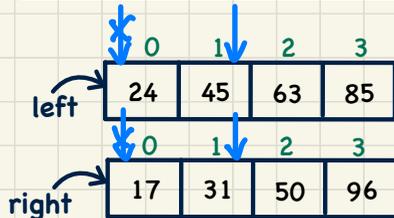
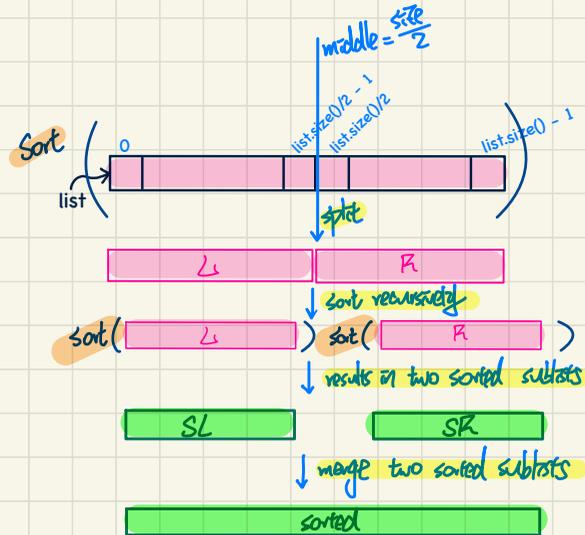
Merge Sort

Merge Sort in Java

```

public List<Integer> sort(List<Integer> list) {
    List<Integer> sortedList;
    if(list.size() == 0) { sortedList = new ArrayList<>(); }
    else if(list.size() == 1) {
        sortedList = new ArrayList<>();
        sortedList.add(list.get(0));
    }
    else {
        int middle = list.size() / 2;
        List<Integer> left = list.subList(0, middle);
        List<Integer> right = list.subList(middle, list.size());
        List<Integer> sortedLeft = sort(left);
        List<Integer> sortedRight = sort(right);
        sortedList = merge(sortedLeft, sortedRight);
    }
    return sortedList;
}
    
```

base cases



Precondition
L and R sorted



```

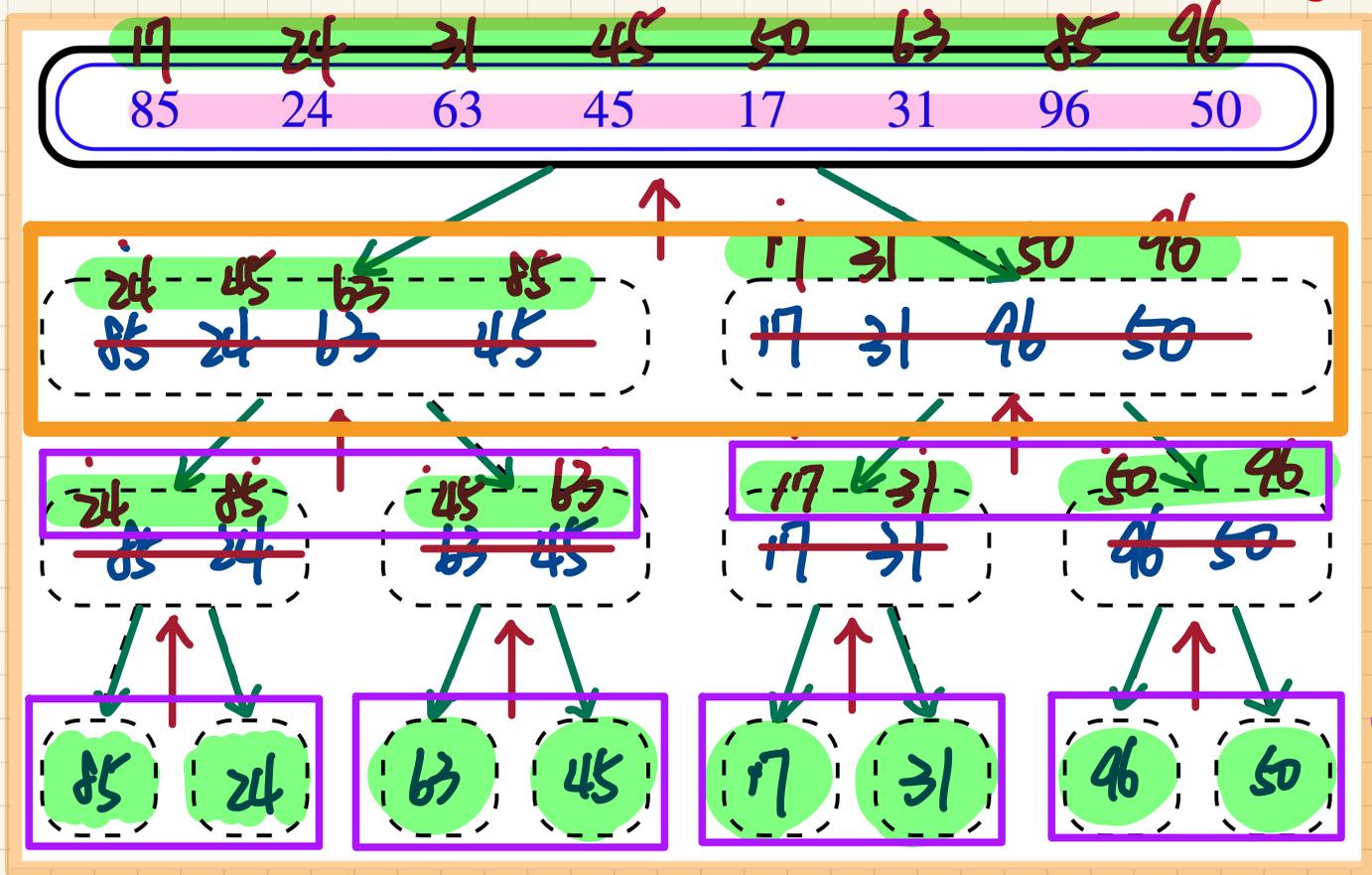
/* Assumption: L and R are both already sorted. */
private List<Integer> merge(List<Integer> L, List<Integer> R) {
    List<Integer> merge = new ArrayList<>();
    if(L.isEmpty() || R.isEmpty()) { merge.addAll(L); merge.addAll(R); }
    else {
        int i = 0;
        int j = 0;
        while(i < L.size() && j < R.size()) {
            if(L.get(i) <= R.get(j)) { merge.add(L.get(i)); i++; }
            else { merge.add(R.get(j)); j++; }
        }
        /* If i >= L.size(), then this for loop is skipped. */
        for(int k = i; k < L.size(); k++) { merge.add(L.get(k)); }
        /* If j >= R.size(), then this for loop is skipped. */
        for(int k = j; k < R.size(); k++) { merge.add(R.get(k)); }
    }
    return merge;
}
    
```

(a) # iterations: $\min(L.size(), R.size())$
 (b) # iterations: remaining # of items to loop over in the longer list.
 $(a) + (b) = L.size() + R.size()$

Merge Sort: Tracing

→ split

→ merge



$\Theta(O(n))$

$\Theta(O(n))$



Merge Sort: Running Time

size = $\frac{n}{2}$

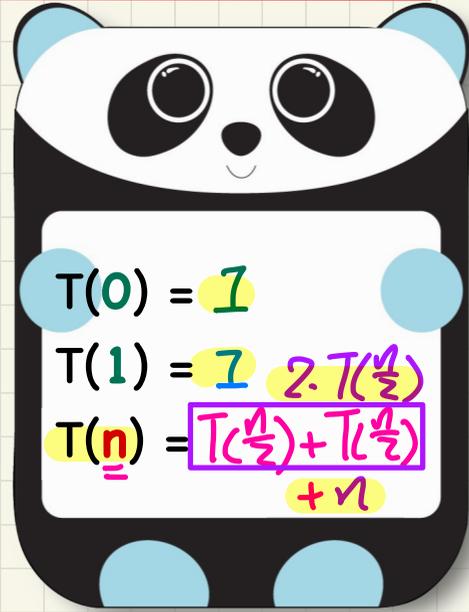
```

public List<Integer> sort(List<Integer> list) {
    List<Integer> sortedList;
    if(list.size() == 0) { sortedList = new ArrayList<>(); }
    else if(list.size() == 1) {
        sortedList = new ArrayList<>();
        sortedList.add(list.get(0));
    }
    else {
        int middle = list.size() / 2;
        List<Integer> left = list.subList(0, middle);
        List<Integer> right = list.subList(middle, list.size());
        List<Integer> sortedLeft = sort(left);
        List<Integer> sortedRight = sort(right);
        sortedList = merge(sortedLeft, sortedRight);
    }
    return sortedList;
}
    
```

sort(left)
sort(right)

recursion tree is a full BT

Running Time as a Recurrence Relation

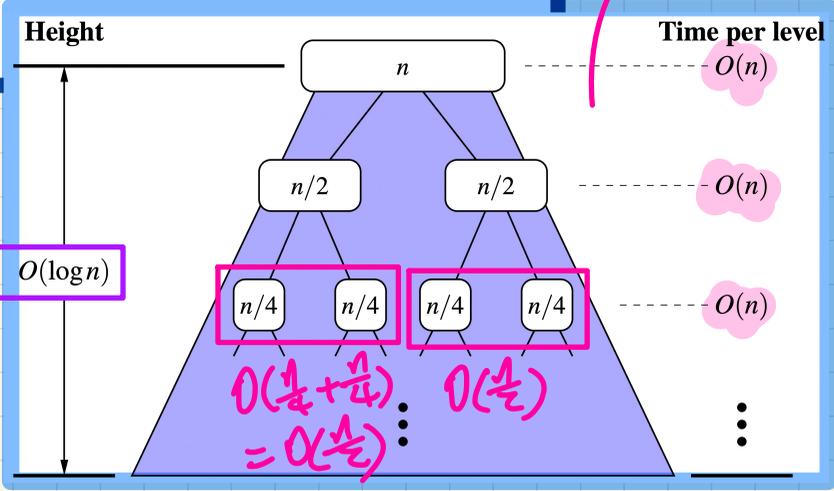


$$T(0) = 1$$

$$T(1) = 1 + 2 \cdot T\left(\frac{n}{2}\right)$$

$$T(n) = T\left(\frac{n}{2}\right) + T\left(\frac{n}{2}\right) + n$$

Total RT:
 $O(\log n \times n)$
 $= O(n \cdot \log n)$
height of balanced BST



Running Time: Unfolding Recurrence Relation

$$T(0) = 1$$

$$T(1) = 1$$

$$T(n) = 2 \cdot T(n/2) + n$$

$$I = \frac{n}{n} = \frac{n}{2^{\log n}}$$

$$n=8 \\ 2^{\log 8} = 8$$

$$T(n) = 2 \cdot T\left(\frac{n}{2}\right) + n$$

$$= 2 \cdot \left(2 \cdot T\left(\frac{n}{4}\right) + \frac{n}{2}\right) + n \quad [4 \cdot T\left(\frac{n}{4}\right) + 2n]$$

$$= 2 \cdot \left(2 \cdot \left(2 \cdot T\left(\frac{n}{8}\right) + \frac{n}{4}\right) + \frac{n}{2}\right) + n \quad [8 \cdot T\left(\frac{n}{8}\right) + 3n]$$

⋮

$$= \frac{2^{\log n}}{n} \cdot T(1) + \log n \cdot n = n + n \cdot \log n \\ = O(n \cdot \log n)$$



Lecture

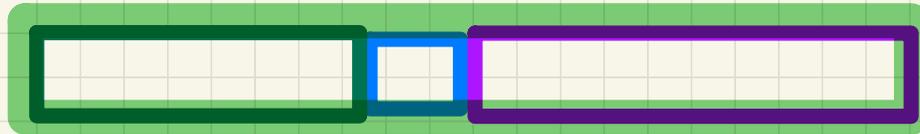
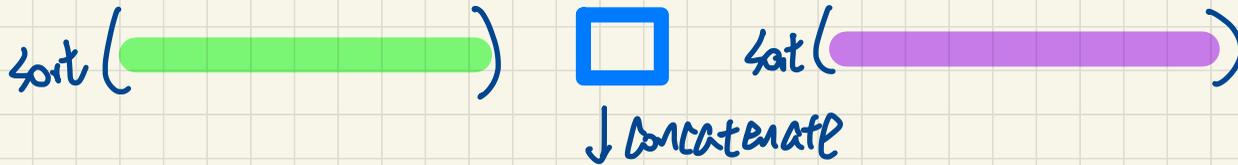
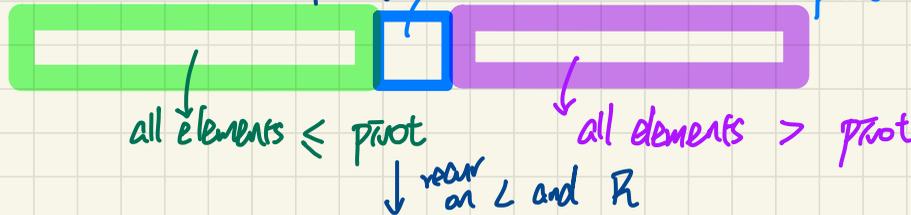
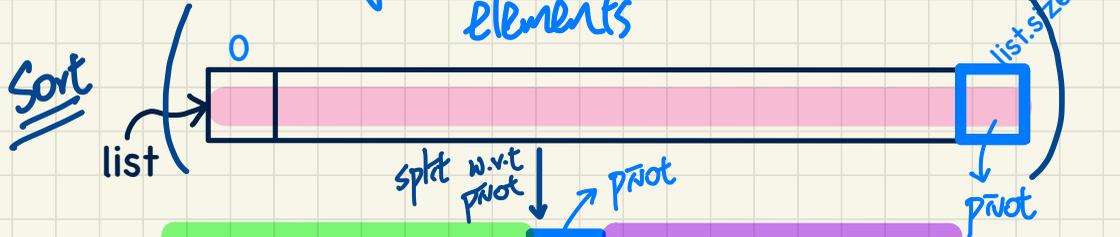
Recursion: Part II (continued)

Quick Sort

Quick Sort: Ideas



pivot: ideally the median value of the list elements



sorted version of input list.

Quick Sort in Java

```

public List<Integer> sort(List<Integer> list) {
    List<Integer> sortedList;
    if(list.size() == 0) { sortedList = new ArrayList<>(); }
    else if(list.size() == 1) {
        sortedList = new ArrayList<>(); sortedList.add(list.get(0));
    }
    else {
        int pivotIndex = list.size() - 1;
        int pivotValue = list.get(pivotIndex);
        List<Integer> left = allLessThanOrEqualTo(pivotIndex, list);
        List<Integer> right = allLargerThan(pivotIndex, list);
        List<Integer> sortedLeft = sort(left);
        List<Integer> sortedRight = sort(right);
        sortedList = new ArrayList<>();
        sortedList.addAll(sortedLeft);
        sortedList.add(pivotValue);
        sortedList.addAll(sortedRight);
    }
    return sortedList;
}
    
```

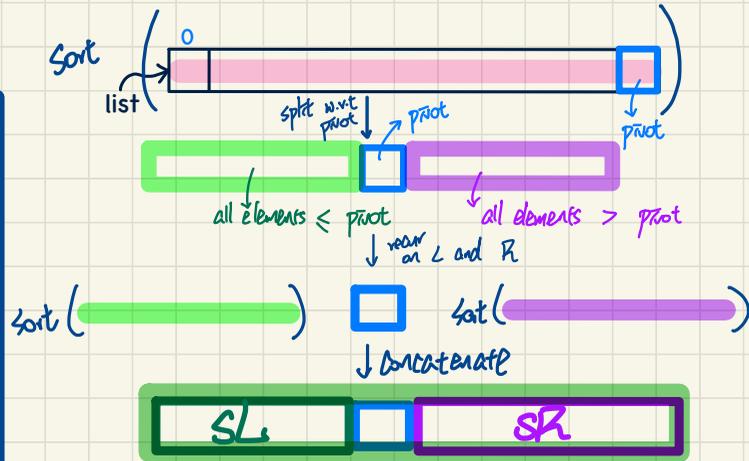
base cases

0(1)

1. Best case: pivot is st. $|L| \approx |R|$

2. Worst case: $|L| \ll |R|$ or $|R| \ll |L|$

OLN



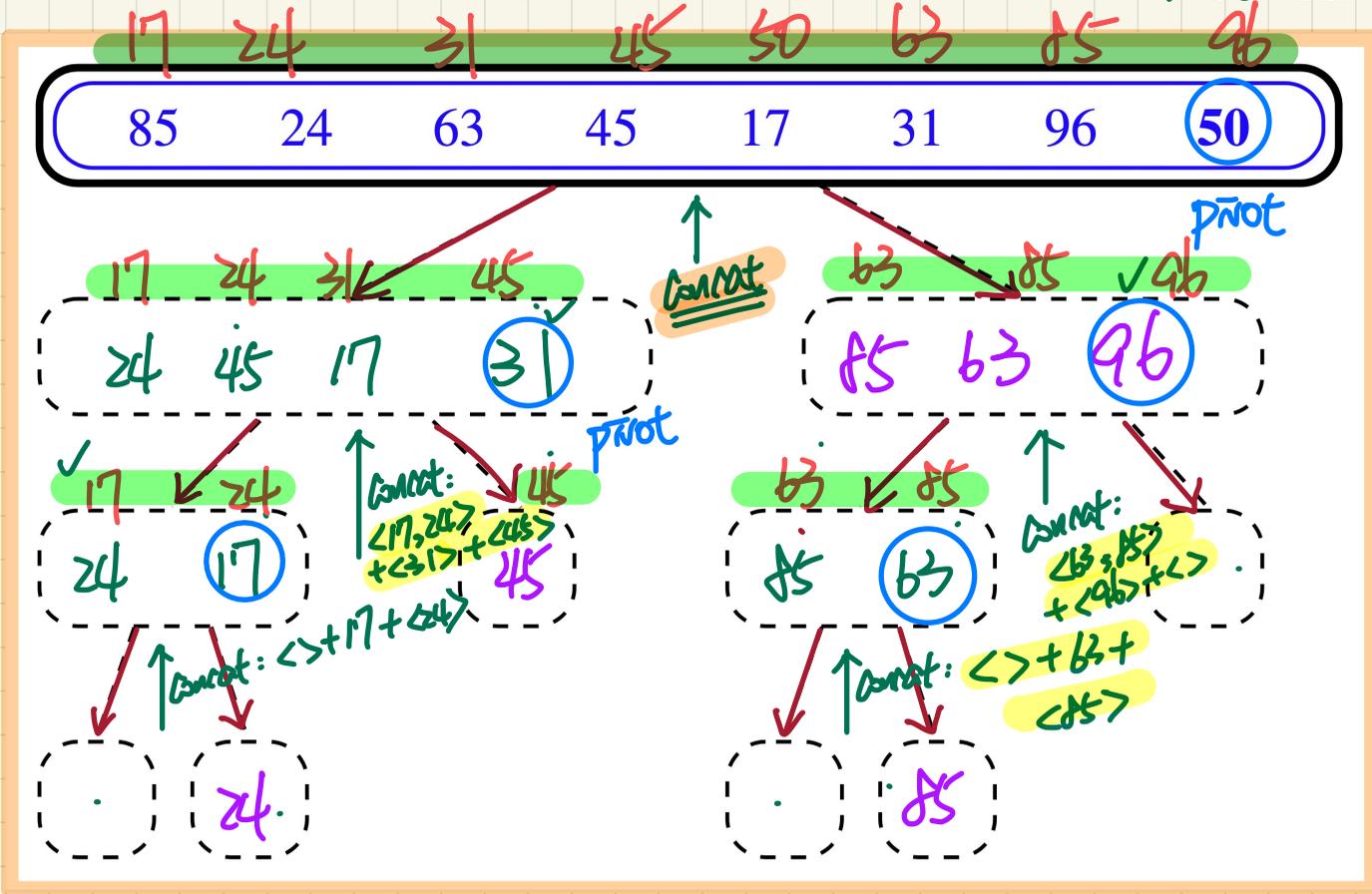
```

List<Integer> allLessThanOrEqualTo(int pivotIndex, List<Integer> list) {
    List<Integer> sublist = new ArrayList<>();
    int pivotValue = list.get(pivotIndex);
    for(int i = 0; i < list.size(); i++) {
        int v = list.get(i);
        if(i != pivotIndex && v <= pivotValue) { sublist.add(v); }
    }
    return sublist;
}

List<Integer> allLargerThan(int pivotIndex, List<Integer> list) {
    List<Integer> sublist = new ArrayList<>();
    int pivotValue = list.get(pivotIndex);
    for(int i = 0; i < list.size(); i++) {
        int v = list.get(i);
        if(i != pivotIndex && v > pivotValue) { sublist.add(v); }
    }
    return sublist;
}
    
```

Quick Sort: Tracing

→ split
→ concatenate



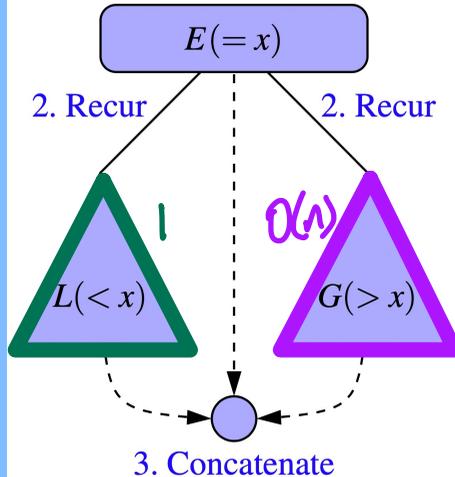
Quick Sort: Worst-Case Running Time

```

public List<Integer> sort(List<Integer> list) {
    List<Integer> sortedList;
    • if(list.size() == 0) { sortedList = new ArrayList<>(); }
    • else if(list.size() == 1) {
        sortedList = new ArrayList<>(); sortedList.add(list.get(0)); }
    else {
        int pivotIndex = list.size() - 1;
        int pivotValue = list.get(pivotIndex);
        List<Integer> left = allLessThanOrEqualTo(pivotIndex, list);
        List<Integer> right = allLargerThan(pivotIndex, list);
        List<Integer> sortedLeft = sort(left);
        List<Integer> sortedRight = sort(right);
        sortedList = new ArrayList<>();
        sortedList.addAll(sortedLeft);
        sortedList.add(pivotValue);
        sortedList.addAll(sortedRight);
    }
    return sortedList;
}
    
```

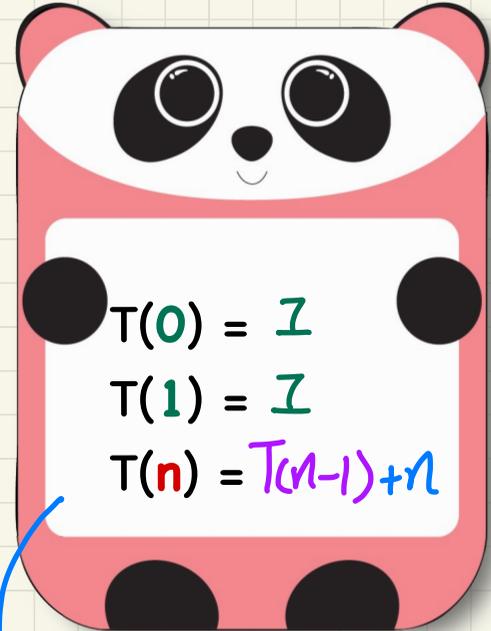
$O(n)$

1. Split using pivot x



splits:
4 $O(n)$

Running Time as a Recurrence Relation



$$T(0) = 1$$

$$T(1) = 1$$

$$T(n) = T(n-1) + n$$

Exercise: Solve by unrolling

Quick Sort: Best-Case Running Time

log₂n

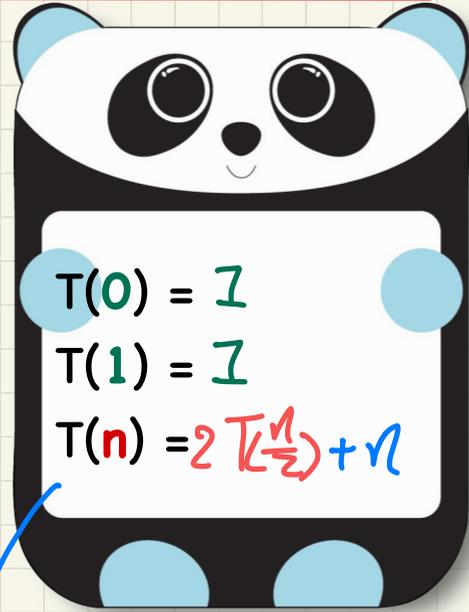
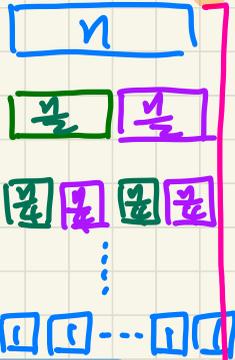
Running Time as a Recurrence Relation

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    List<Integer> sortedList;
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        sortedList = new ArrayList<>(); sortedList.add(list.get(0));
    }
    else {
        int pivotIndex = list.size() - 1;
        int pivotValue = list.get(pivotIndex);
        List<Integer> left = allLessThanOrEqualTo(pivotIndex, list);
        List<Integer> right = allLargerThan(pivotIndex, list);
        List<Integer> sortedLeft = sort(left);
        List<Integer> sortedRight = sort(right);
        sortedList = new ArrayList<>();
        sortedList.addAll(sortedLeft);
        sortedList.add(pivotValue);
        sortedList.addAll(sortedRight);
    }
    return sortedList;
}
    
```

$O(1)$

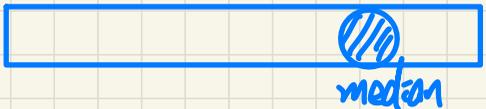
$O(n)$



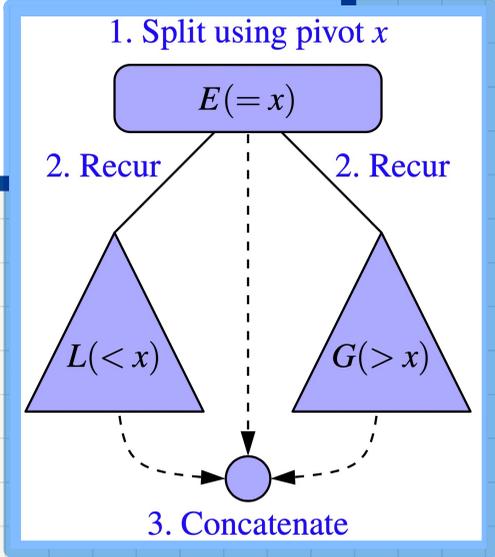
$$T(0) = 1$$

$$T(1) = 1$$

$$T(n) = 2T\left(\frac{n}{2}\right) + n$$



sizes equal



Ex. 2 Exercise: solve by Unfolding.

$$T(0) = 1$$

$$T(1) = 1$$

$$T(n) = T\left(\frac{n}{2}\right) + T\left(\frac{n}{2}\right) + n$$